

### Problem Set 3 – Due Monday, October 13, 2008

1. (Velleman, p. 64, problems 8 and 9.) Complete the following table, answering whether the statement is true (True) or false (False) when the universe of discourse is as indicated.

	$\mathbb{R}$	$\mathbb{Z}$
$\forall x \exists y (2x - y = 0)$		
$\exists y \forall x (2x - y = 0)$		
$\forall x \exists y (x - 2y = 0)$		
$\forall x (x < 10 \rightarrow \forall y (y < x \rightarrow y < 9))$		
$\exists y \exists z (y + z = 100)$		
$\forall x \exists y (y > x \wedge \exists z (y + z = 100))$		

2. (Velleman, p. 72, problem 2.) Translate the negation of the following statements into formulas of quantification logic, introducing predicates as needed.
- There is someone in the freshman class who doesn't have a roommate.
  - Everyone likes someone, but no one likes everyone.
  - $(\forall a \in A)(\exists b \in B)(a \in C \leftrightarrow b \in C)$
  - $(\forall y > 0)(\exists x)(ax^2 + bx + c = y)$
3. Translate the following into a formula of the predicate calculus. "A language  $L$  that is regular will have the following property: there will be some number  $N$  (that depends on  $L$ ) such that if  $s$  is a string in  $L$  (a *string* is a sequence of characters) whose length is at least  $N$  then  $s$  can be written as  $xyz$  where  $y$  is not the empty string and  $xy^iz$  is in the language  $L$  for every nonnegative integer  $i$ ." *Note 1: you do not have to understand this statement to do this problem. Note 2: you will see this specific statement again in ECS 120.*
4. A well-formed formula is said to be in *conjunctive normal form* (CNF) if it is the conjunct (and) of terms where each term is the disjunct (or) of variables or their complements. Convert the following formula into CNF:  $\phi = A \wedge (B \rightarrow C)$ . Can every formula be converted into a logically equivalent one in CNF? Explain your answer.
5. In class we described the problem TILING, where you are given a collection of  $t$  tile types,  $1, \dots, t$ , and Boolean functions  $S(k, k')$  and  $T(k, k')$  with the following semantics:  $S(k, k')$  is true iff a tile of type  $k$  may be put immediately to the left of a tile of type  $k'$ , and  $T(k, k')$  is true iff a tile of top  $k$  may be put immediately beneath a tile of type  $k'$ . Explicitly specify a collection of formulas  $\Gamma$  such that  $\Gamma$  is satisfiable (meaning that there is a truth assignment that satisfies every formula in  $\Gamma$ ) iff the upper quadrant of the plane is tilable according to the  $S$  and  $T$  constraints.