

## Midterm — Section 1

**Instructions:** Please answer the questions succinctly and thoughtfully. Good luck.

— Phil Rogaway

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**Name:**

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**Signature:**

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On problem	you got	out of
1		45
2		30
3		25
$\Sigma$		100

**1 Short Answer****[45 points]**

- (1) Let  $M_1$  be an  $n_1$ -state DFA and let  $M_2$  be an  $n_2$ -state DFA. Using the procedures given in class and in your text, how many states will be in the NFA  $M$  the language of which is  $L(M_1) \cup (L(M_2))^*$ ? Do not make your machine unnecessarily large. Explain your reasoning.

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- (2) Using the procedure we have seen in class, convert the regular expression  $1^* \cup 00$  into an NFA for the same language.

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- (3) Complete the following, mathematically precise, definition, according to the conventions of our text:

an **NFA** is a 5-tuple  $M = ($  ) where:

(4) Recall that, for  $L \subseteq \{0, 1\}^*$ ,  $\text{PAL}(L) = \{x \in \{0, 1\}^* : xx^R \in L\}$ . Write a regular expression for  $\text{PAL}(\{\varepsilon, 00, 01, 000, 0110\})$ .

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(5) Carefully describe a decision procedure (algorithm) to decide the following language:  $L = \{\langle \alpha \rangle : \alpha \text{ is a regular expression and the smallest NFA for } L(\alpha) \text{ is smaller than the smallest DFA for } L(\alpha)\}$ . (smallest = fewest states)

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(6) Write a CFG for the language  $L = \{1^i \neq 1^j : i, j \geq 1, i \neq j\}$ . Use no more than 6 rules.

- (7) Use the pumping lemma for regular languages to prove that the following language is **not** regular:  $L = \{1^i < 1^j : 1 \leq i < j\}$ .

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- (8) Prove or disprove: the following grammar is **ambiguous**:

$$S \rightarrow S1S1S2S \mid S1S2S1S \mid S2S1S1S \mid \varepsilon$$

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- (9) Carefully **state** the pumping lemma for context-free languages.

**2 Justified True or False** [30 points]

Put an **X** through the **correct** box. Where it says “Explain” provide a **brief** (but convincing) justification. No credit will be given to correct answers that lack a proper justification. Where appropriate, **make your justification a counter-example**. Throughout, we use  $L$  to denote a language (maybe regular, maybe not).

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1. Every CFL has some ambiguous grammar.

**True****False**

Explain:

2. Every subset of a regular language is regular.

**True****False**

Explain:

3. For every  $n \geq 1$ , there is an  $2n$ -state DFA for the language  $0\{0, 1\}^n$ .

Explain:

**True****False**

4. The intersection of two context-free languages is context free.

**True****False**

Explain:

5. In class we exhibited an algorithm to decide if two context-free grammars generate the same language.

**True**

**False**

6. If  $L$  is context free then some NPDA accepts the complement of  $L$ .

**True**

**False**

**3 A Little Proof****[25 points]**

Find the smallest DFA for the language  $L = (000)^* \cup (111)^*$ . Prove that there is no DFA smaller than yours. (smaller=fewer states.)