CAV23

# Incremental Dead State Detection in Logarithmic Time

# Caleb Stanford and Margus Veanes





#### Dead State Detection in Automata



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## **Existing Solutions**



Best result in online graph algorithms:O(sqrt m) per edge[Bender, Fineman, Gilbert, Tarjan 2015]

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#### Our main result: O(log m) per edge

# Talk Outline

- 1. Motivation
- 2. Guided Incremental Digraphs (GID)
  - Broadly applicable data structure for dead state detection
- 3. Algorithms
- 4. Evaluation
  - **110-530x speedup** over [BFGT 2015]

#### [PLDI 2021]



#### Motivation in Z3

#### 2020 internship







Boolean regex constraints: s matches R1 and R2 and ... And does not match R4, R5, ...

#### [PLDI 2021]



#### Motivation in Z3

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Boolean regex constraints: s matches R1 and R2 and ... And does not match R4, R5, ...



"The total number of invocations of Zelkova ranges from a few million to tens of millions in a single day"



#### [PLDI 2021]

## Motivation in Z3

# **Z**3

# And does not match R4, R5, ... **Exponential blowup Too expensive!**

s matches R1 and R2 and ...





#### **Derivatives:**







#### **Derivatives:**

23





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R = aΣ | b((ΣΣ)\*  $\cap$  Σ(ΣΣ)\*)





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 $\mathsf{R} = \mathsf{a}\Sigma \mid \mathsf{b}((\Sigma\Sigma)^* \cap \Sigma(\Sigma\Sigma)^*)$ 



*Lazy* decision procedure – very fast in practice!



[PLDI 2021]

#### **Derivatives:**

 $\mathsf{R} = \mathsf{a}\Sigma \mid \mathsf{b}((\Sigma\Sigma)^* \cap \Sigma(\Sigma\Sigma)^*)$ 



*Lazy* decision procedure – very fast in practice!

How do we detect dead states?







ly

# Maintaining a topological order under edge insertions

Alberto Marchetti-Space

A New Approach to Incremental Cycle Detection and Related Problems

#### Incremental Cycle Detectio Component Maintenance

BERNHARD HAEUPLER, Massachi TELIKEPALLI KAVITHA, Tata Institute ROGERS MATHEW, Indian Institute SIDDHARTHA SEN, Princeton Unive ROBERT E. TARJAN, Princeton Unive

Michael A. Bender Department of Computer Science Stony Brook University Jeremy T. Fineman Department of Computer Science Georgetown University

Seth Gilbert Department of Computer Science National University of Singapore

Robert E. Tarjan HP and Department of Computer Science Princeton University



#### **Dead State Detection**



#### **Dead State Detection**



Basically a naive BFS/DFS O(m) per edge

Fast forward 3 years...

# **Defining the Problem**

Almost every problem that you come across is befuddled with all kinds of extraneous data of one sort or another; and if you can bring this problem down into the main issues, you can see more clearly what you're trying to do.

-Claude Shannon

# Simplifying...



# Simplifying...



# Simplifying...



#### Guided Incremental Digraph (GID)

Closed Open

Outgoing edges not allowed!

## Guided Incremental Digraph (GID)



Live: can reach a terminal state

Dead: not live and all reachable states are closed

Closed

Open

## Solving the Problem



#### Solving the Problem

BFS/DFS: O(m) per update



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BFS/DFS: O(m) per update

Maintain the graph as a set of SCCs [BFGT2015]

- O(sqrt(m)) per update



What information do we need for non-dead states?



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Directed rooted forest









#### O(m) to check for cycle



Clever reduction to **undirected reachability** for undirected forests

#### Are we done?

Asymptotic complexity: O(log m) amortized per graph update

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- ...and they impose data structure overheads

#### Euler Tour Trees 1510

LoC

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#### Solution: A second, *lazy* algorithm – efficient in practice

# **Evaluation**

## Evaluation

How does it perform compared to state-of-the-art online graph algorithms?

In the paper:

- How does performance change with the graph class?
- How does it perform on graphs from the Z3 regex application?



**Benchmark Size** 

#### High-level takeaways

- Online graph algorithms are useful in formal methods
- Incremental dead state detection is a natural problem that arises in

practical verification tools

• Asymptotic complexity is not always enough

# Summary

#### **Guided Incremental Digraphs**

- Closed states: no more outgoing edges

New algorithms: for dead state detection

- In log(m) time
- + practical improvements

Publicly available on GitHub and crates.io

- https://github.com/cdstanford/gid



#### Future Work

Does this generalize to other problems like minimization?

Lazy decision procedures for other contexts

(e.g., LTL and Büchi automata)



Results

