

CS 224 Fall 2009, HW 4 Due November 5, although that might not be enough time - let me know.

## Splits Equivalence

Problems 1, 2: In the notes on the Perfect Phylogeny Problem, on pages 18 and 19 there are two problems concerning Splits Equivalence and its relation to the Four Gametes Theorem. Solve those problems.

## Steiner Tree problem on Hypercubes

Let  $G = (N, E)$  be an undirected graph on node set  $N$ . Let  $X \subseteq N$  be a given subset of nodes. A *Steiner tree*  $ST$  of  $G$  for  $X$  is any connected subtree of  $G$  that contains all the nodes of  $X$ , although it may contain nodes in  $N - X$  as well. The *weight* of a Steiner tree  $ST$  is the number of edges in  $ST$ , and is denoted  $W(ST)$ . The *Steiner tree problem*, given  $G$  and  $X$ , is to find the Steiner tree for  $X$  in  $G$  of minimum weight.

A *hypercube* of dimension  $k$  is an undirected graph with  $2^k$  nodes, where the nodes are labeled with the integers between 0 and  $2^k - 1$ . Two nodes in the hypercube are adjacent if and only if the binary representation of their labels differs in exactly one bit. The Steiner tree problem on hypercubes is the Steiner tree problem where the graphs are hypercubes.

The (two-state) perfect phylogeny problem can be viewed as a question about the weight of minimum weight Steiner Tree.

Problem 3: What question about the minimum weight Steiner Tree can be cast as the perfect phylogeny problem? As a hint, develop a simple lower bound on  $W(ST)$  that applies to any  $G$  and  $X$ . Then relate the perfect phylogeny problem to that lower bound.

## Rooted Tree Compatibility

**Definition** A rooted, leaf-labeled tree  $T'$  is a *refinement* of another rooted, leaf-labeled tree  $T$  if  $T$  can be obtained by a series of contractions of edges of  $T'$ . When an edge  $(u, v)$  touching a leaf  $v$  is contracted, the label of  $v$  is

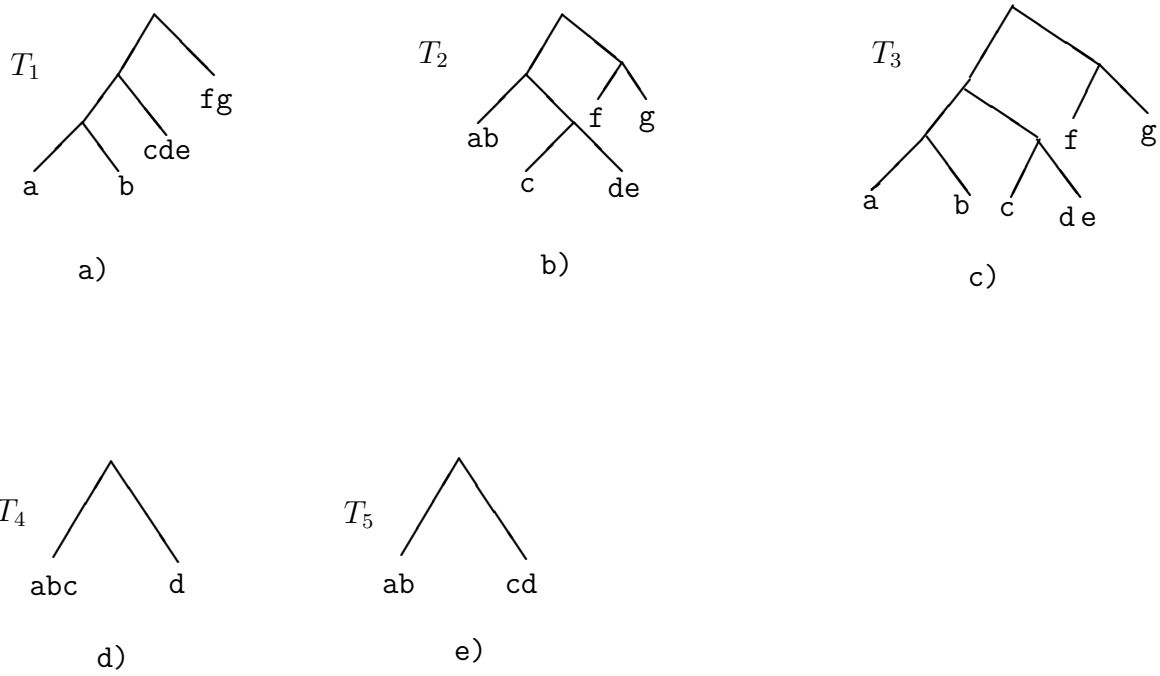


Figure 1: Trees  $T_1$  and  $T_2$  are compatible; they are refined by  $T_3$ . Trees  $T_4$  and  $T_5$  are not compatible.

added to any label on  $u$ . In general when an edge  $(u, v)$  directed from  $u$  to  $v$  is contracted, any label on  $v$  is added to whatever label is on  $u$ .

If  $T'$  refines  $T$ , then  $T'$  agrees with all the evolutionary history displayed in  $T$ , while displaying additional history not contained in  $T$ .

Let  $T_1$  and  $T_2$  be two rooted trees whose leaves are labeled with the same set of labels. But note that  $T_1$  and  $T_2$  can have different numbers of leaves, and a leaf can have more than a single label. We will assume that  $T_1$  and  $T_2$  are both in “reduced form”, that is, both are binary trees, and no node except the root can have exactly one child.

**Definition** Trees  $T_1$  and  $T_2$  are *compatible* if there exists a rooted tree  $T_3$  refining both  $T_1$  and  $T_2$  (see Figure 1).

**Tree compatibility problem** Given trees rooted, leaf-labeled trees  $T_1$  and  $T_2$  on the same label set, determine whether the two trees are compatible, and if so, produce a refinement tree  $T_3$ .

Let  $M_1$  be a 0-1 matrix with one row for each taxon and one column for each node  $j$  in  $T_1$ , including leaves. Entry  $(i, j)$  of  $M_1$  has value one if and only if taxon  $i$  is found at or below node  $j$ . That is, column  $j$  of  $M_1$  records the taxa found in the subtree of  $T_1$  rooted at node  $j$ . Matrix  $M_2$  is similarly defined for  $T_2$ , and matrix  $M_3$  is the matrix formed by the union of the columns of  $M_1$  and  $M_2$ . Then

**Theorem**  $T_1$  and  $T_2$  are compatible if and only if there is a perfect-phylogeny for  $M_3$  with the all-zero ancestral sequence. Further, a perfect-phylogeny  $T_3$  for  $M_3$  is a refinement of both  $T_1$  and  $T_2$ . Note that two columns in  $M_3$  might be identical, and hence  $T_3$  might have a leaf with more than a single label.

Problem 4: Prove (with a good explanation) this theorem.