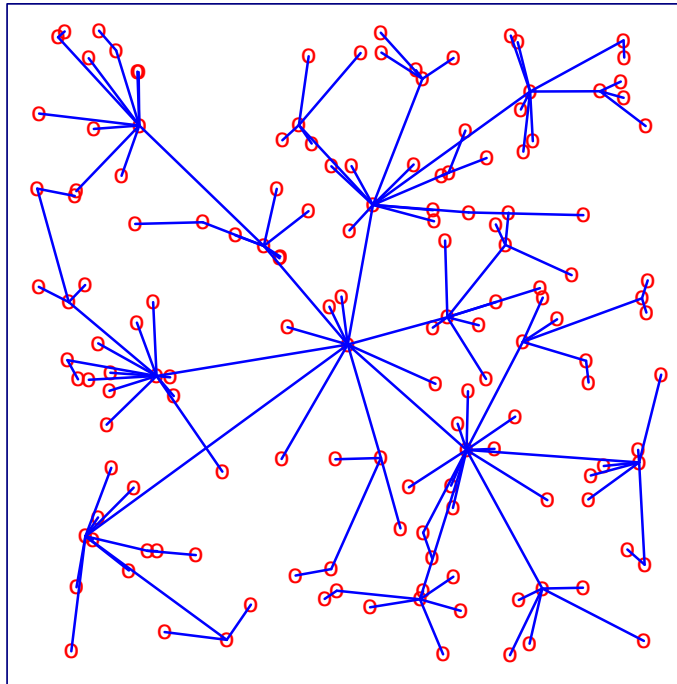


ECS 253 / MAE 253, Lecture 2

April 5, 2023



“Power laws, Random graphs, phase transitions”

Class structure

- Two tracks to the class:

Track A: Project

- (1) Common homeworks (e.g. HW1.pdf, HW2.pdf) and
- (2) HW1a.pdf, HW2a.pdf etc.

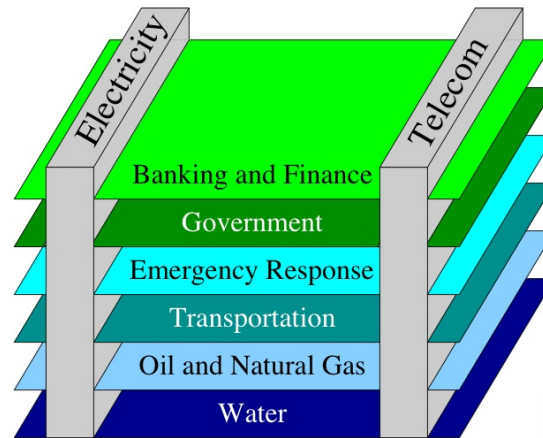
Track B: Advanced HWs

- (1) Common homeworks (e.g. HW1.pdf, HW2.pdf) and
- (2) HW1b.pdf, HW2b.pdf etc.

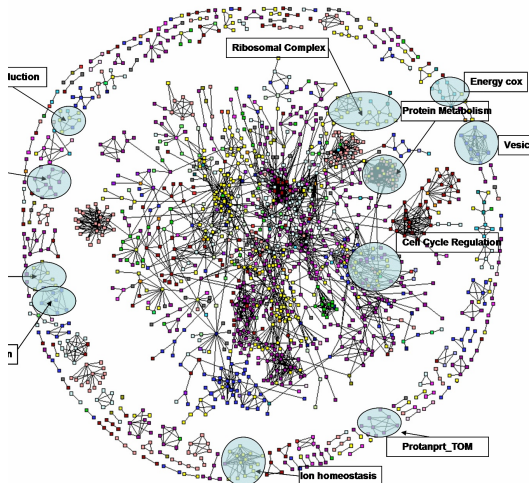
- Track A: Project

- Teams of 5-6 people ideal
- Negative results are OK
- Ideally aim to have a result for a journal or conference

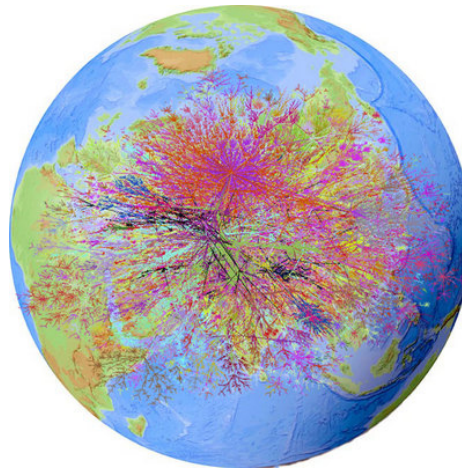
Complex networks are ubiquitous:



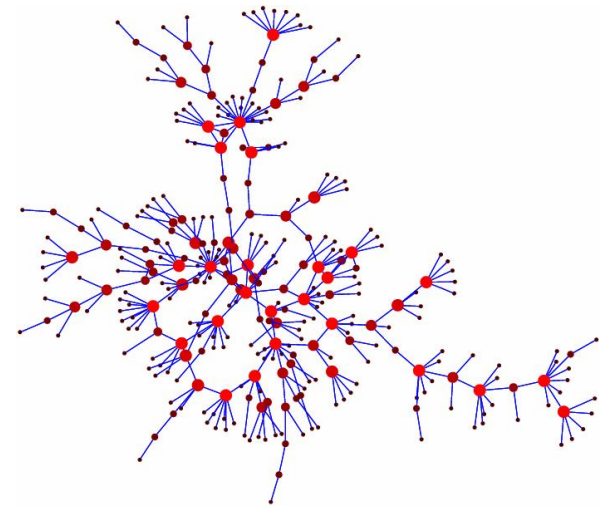
Critical Infrastructure



Biological & Ecological networks



Information and Communication technology



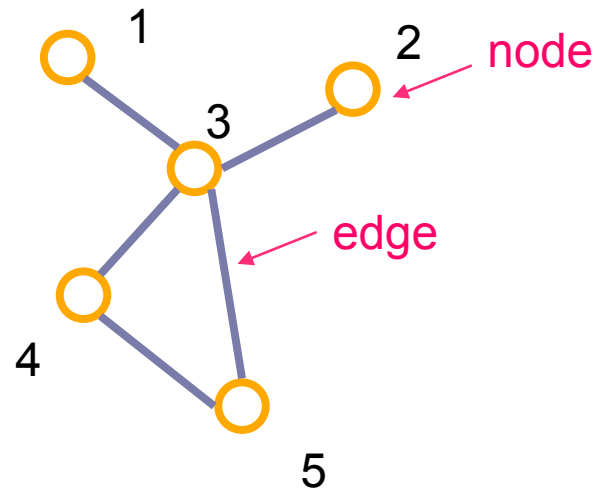
**Social networks:
Economics & Epidemics**

Networks: Physical, Biological, Social, Technological

- **Geometric** versus **virtual** (Internet versus WWW).
- **Natural** /spontaneously arising versus **engineered** /built.
- **Directed** versus **undirected** edges.
- Each network may **optimize** something unique.
- Identifying **similarities** and fundamental **differences** can guide future design/understanding.
- Interplay of **topology** and **function** ?
- Unifying features: – **Broad heterogeneity in node degree.**
– **Small Worlds** (Diameter $\sim \log(N)$).

What are networks?

- Networks are collections of points joined by lines.



“Network” \equiv “Graph”

points	lines	
vertices	edges, arcs	math
nodes	links	computer science
sites	bonds	physics
actors	ties, relations	sociology

Subtle details of edges

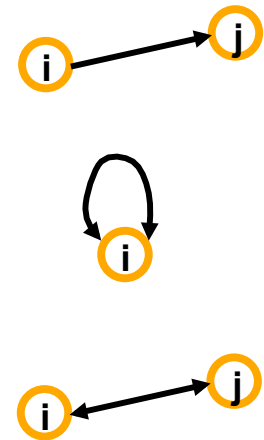
Network elements: edges

- Directed (also called arcs)
 - $A \rightarrow B$ (E_{BA})
 - A likes B, A gave a gift to B, A is B's child
- Undirected
 - $A \leftrightarrow B$ or $A - B$
 - A and B like each other
 - A and B are siblings
 - A and B are co-authors
- Edge attributes
 - weight (e.g. frequency of communication)
 - ranking (best friend, second best friend...)
 - type (friend, relative, co-worker)
 - properties depending on the structure of the rest of the graph:
e.g. betweenness
- Multiedge: multiple edges between two pair of nodes
- Self-edge: from a node to itself

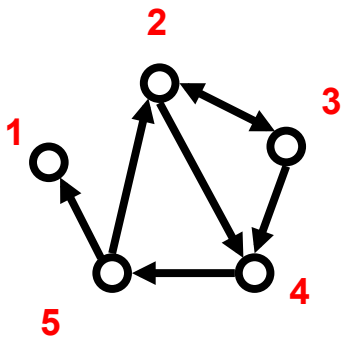
Adjacency matrices

- Representing edges (who is adjacent to whom) as a matrix

- $A_{ij} = 1$ if node i has an edge to node j
= 0 if node i does not have an edge to j
- $A_{ii} = 0$ unless the network has self-loops
 - If self-loop, $A_{ii}=1$
- $A_{ij} = A_{ji}$ if the network is undirected, or if i and j share a reciprocated edge



Example:

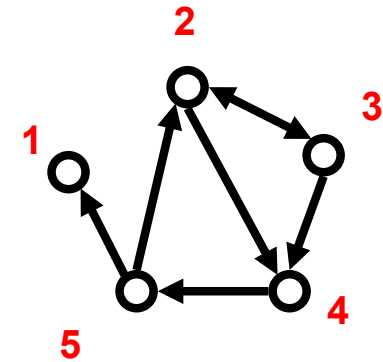


$$A = \begin{pmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 & 0 \end{pmatrix}$$

Adjacency lists

■ Edge list

- 2 3
- 2 4
- 3 2
- 3 4
- 4 5
- 5 2
- 5 1



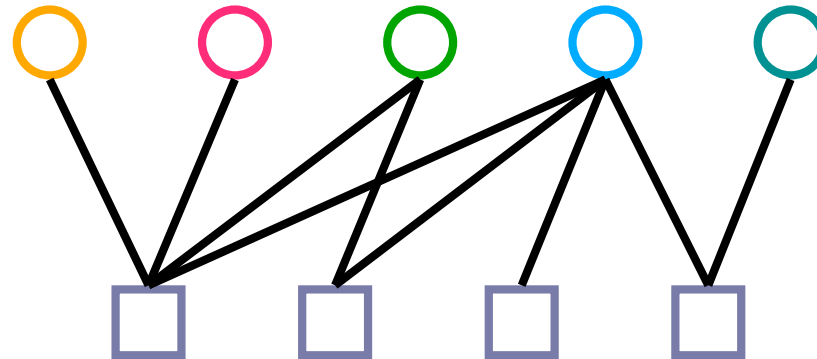
■ Adjacency list

- is easier to work with if network is
 - large
 - sparse
- quickly retrieve all neighbors for a node
 - 1:
 - 2: 3 4
 - 3: 2 4
 - 4: 5
 - 5: 1 2

Beyond simple networks:

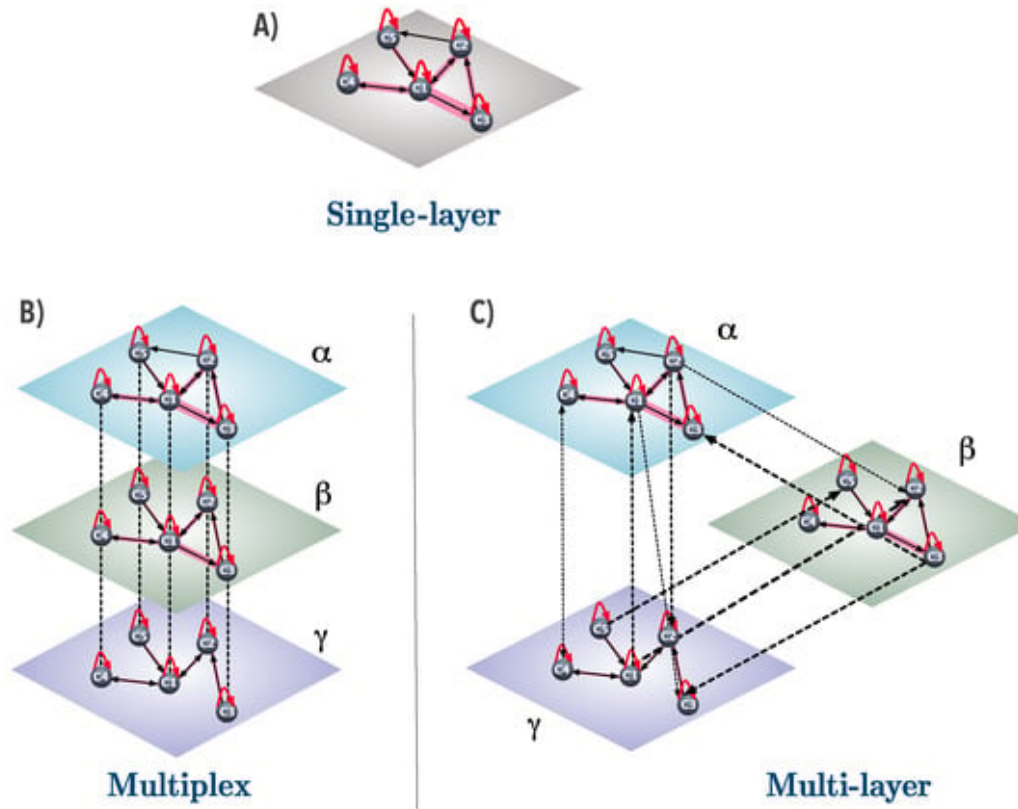
Bipartite (two-mode) networks

- edges occur only between two groups of nodes, not within those groups
- for example, we may have individuals and *events*
 - directors and boards of directors
 - customers and the items they purchase
 - metabolites and the reactions they participate in



Beyond simple networks: Multiplex and multi-layered

Image from Alves et al, Entropy, 2018



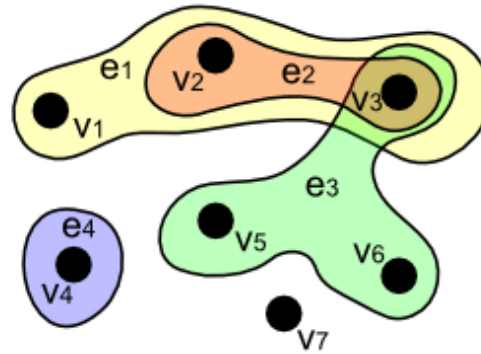
- B) **Multiplex:** the same set of nodes have multiple types of relationships, each one described by a layer.
- C) **Multi-layer:** The nodes in each layer can be distinct.

Beyond simple networks:

HyperGraphs

- Edges join more than two nodes at a time (*hyperEdge*)

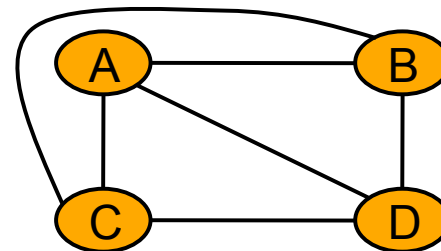
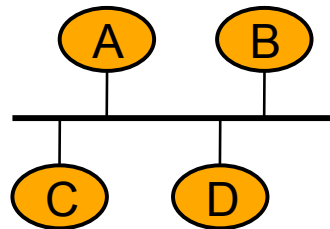
- Affiliation networks



- Examples

- Families
- Subnetworks

- Chemical reactions
- co-author networks

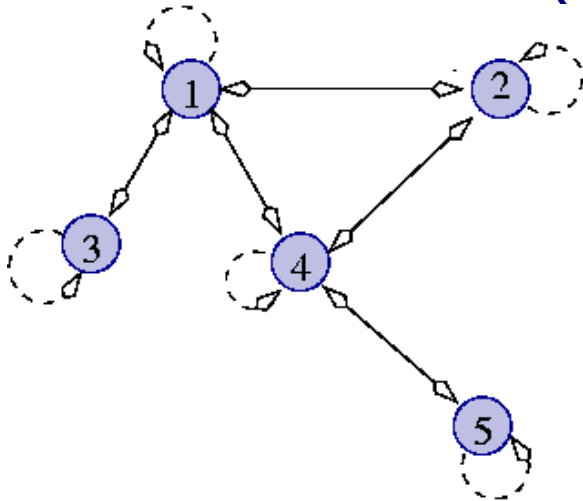


Can be transformed to a *bipartite network*

NETWORK TOPOLOGY; simple edges

Binary connectivity matrix, M :

$$M_{ij} = \begin{cases} 1 & \text{if edge exists between } i \text{ and } j \\ 0 & \text{otherwise.} \end{cases}$$



$$\begin{pmatrix} 1 & 1 & 1 & 1 & 0 \\ 1 & 1 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 & 0 \\ 1 & 1 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 \end{pmatrix} = M$$

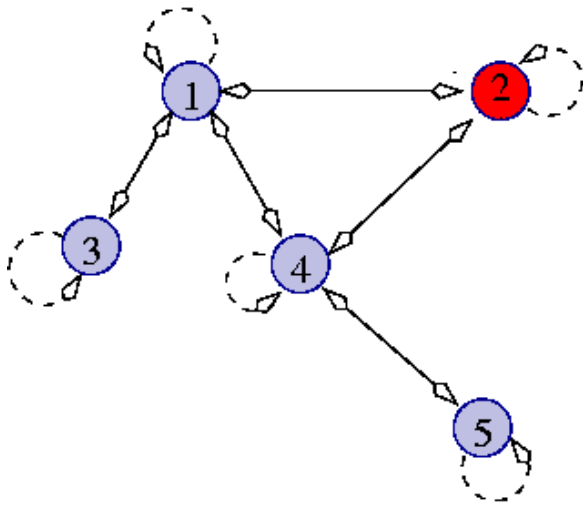
Node **degree** is number of links.

Network Activity: FLOWS on NETWORKS

(Spread of disease, routing data, materials transport/flow, gossip spread/marketing)

FLOWS on NETWORKS : Random walks

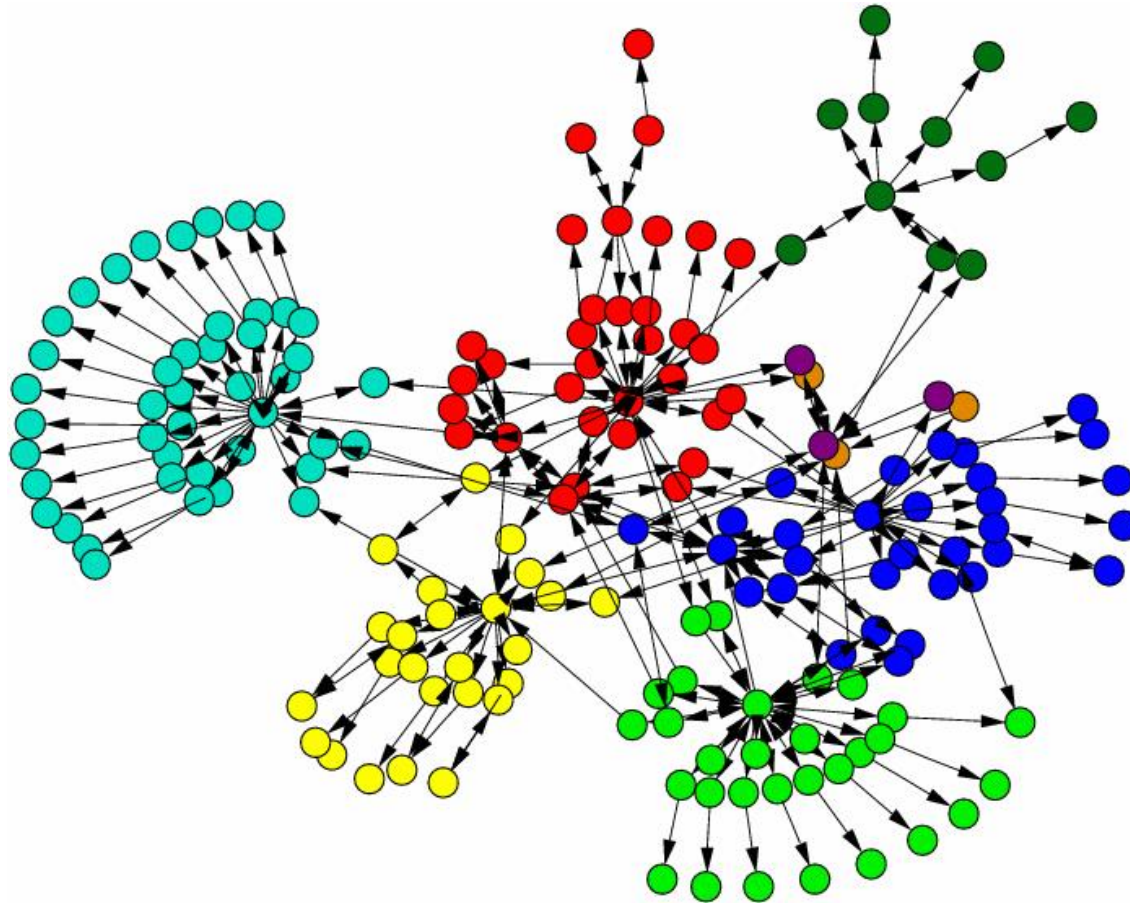
Random walk on the network has state transition matrix, P :
(Column normalize the adjacency matrix)



$$\begin{pmatrix} 1/4 & 1/3 & 1/2 & 1/4 & 0 \\ 1/4 & 1/3 & 0 & 1/4 & 0 \\ 1/4 & 0 & 1/2 & 0 & 0 \\ 1/4 & 1/3 & 0 & 1/4 & 1/2 \\ 0 & 0 & 0 & 1/4 & 1/2 \end{pmatrix} = P$$

The eigenvalues and eigenvectors convey much information.
Markov Chains, Spectral Gap.

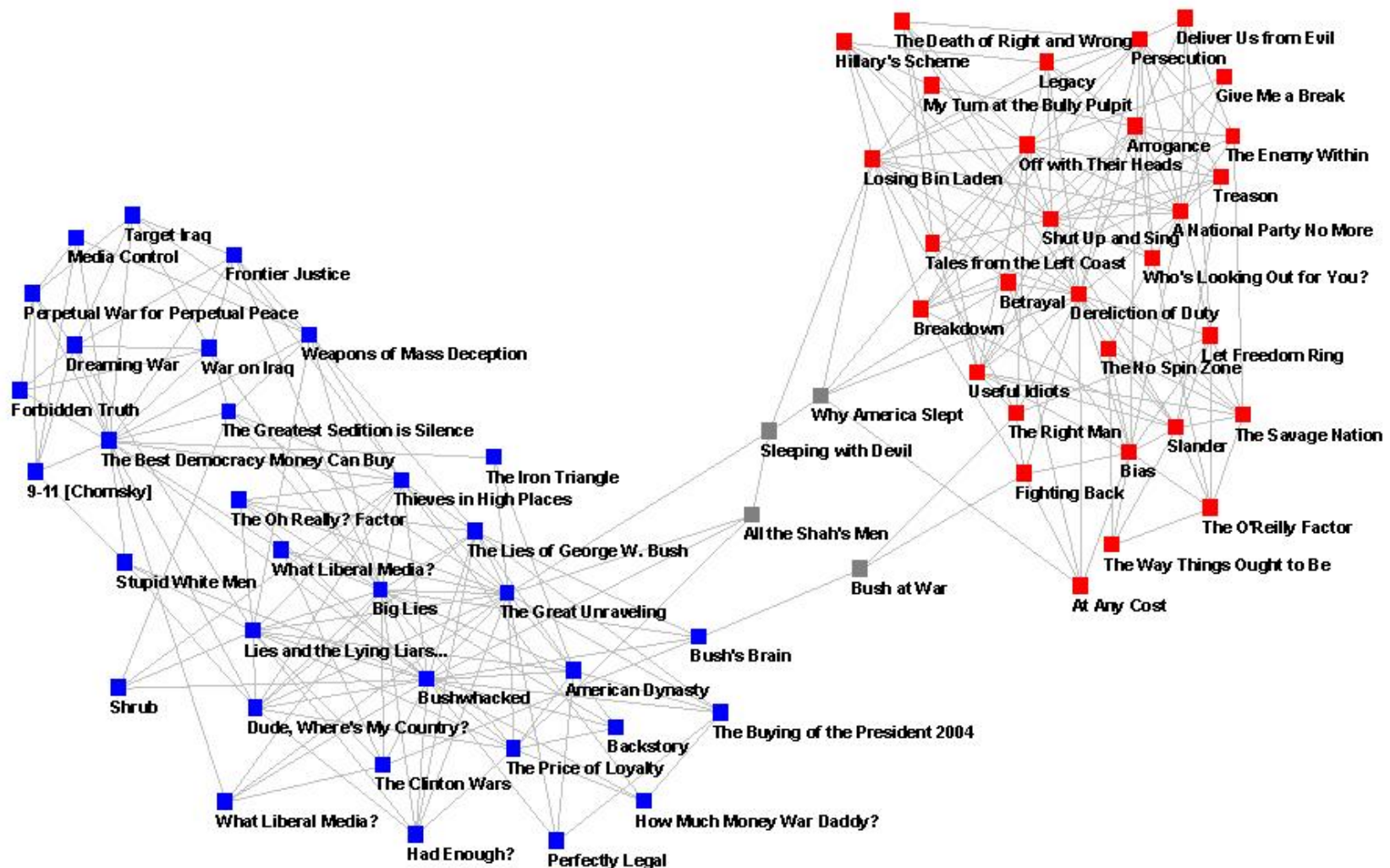
Random walk on the WWW is the “Page Rank”



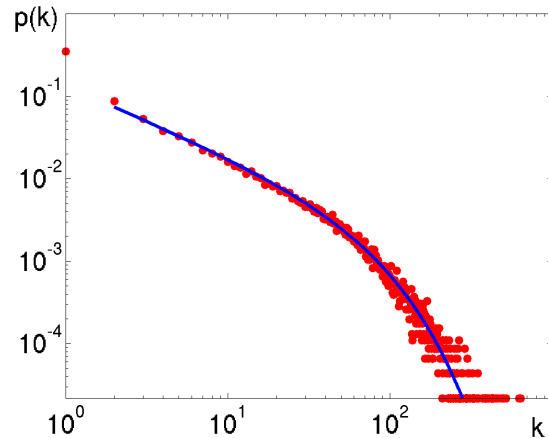
Page Rank of a node is the steady-state random walk occupancy probability.

(We will discuss building a search engine in detail later.)

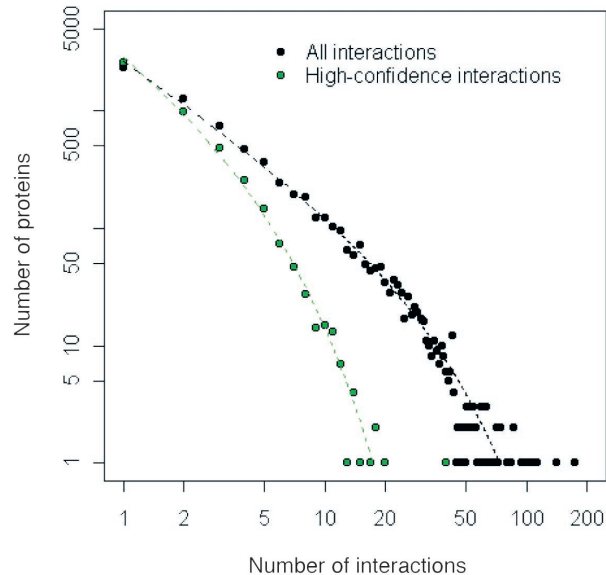
Example Eigen-technique: Community structure (Political Books 2004)



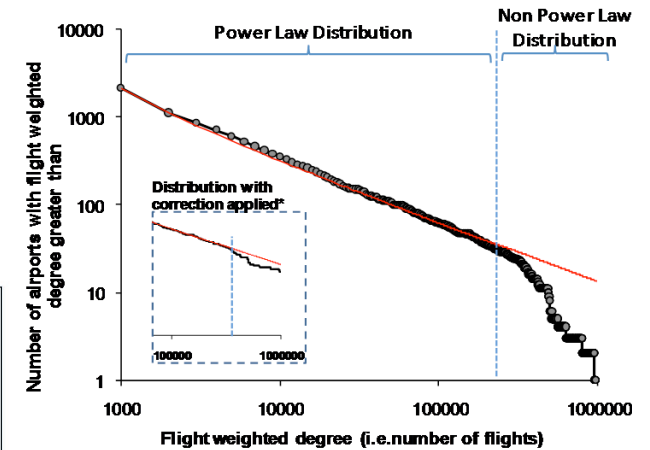
Back to topology: Broad scale degree distributions



Social contacts
Szendrői and Csányi



Protein interactions
Giot et al Science 2003



Airport traffic
Bounova 2009

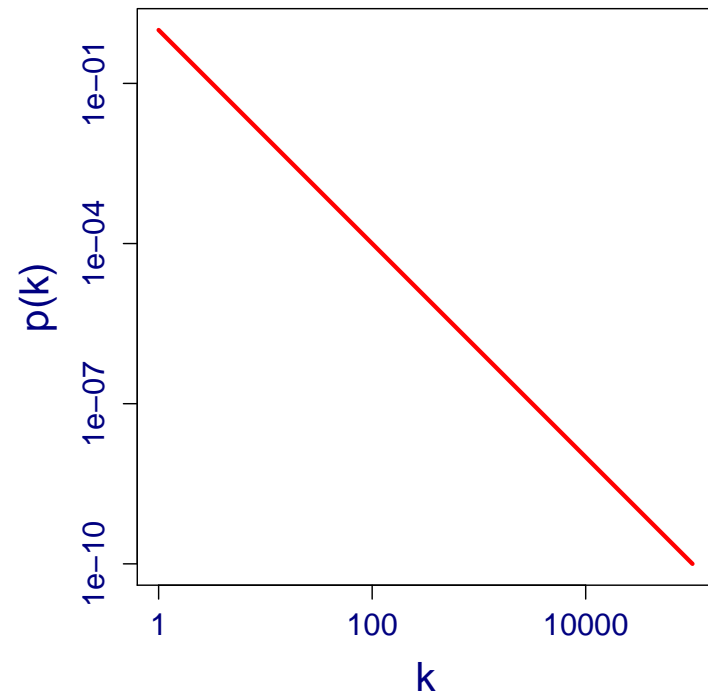
- A few hubs, dominated by leaves
- Small data sets, power laws vs log normal, stretched-exponential, etc...
 - Exceptions: Power grids? Router-level Internet?

Degree distribution

- Often observe “**heavy-tailed**” / “**broad-scale**” degree distributions.
- The simplest example of such a distribution is a power law (Pareto distribution).

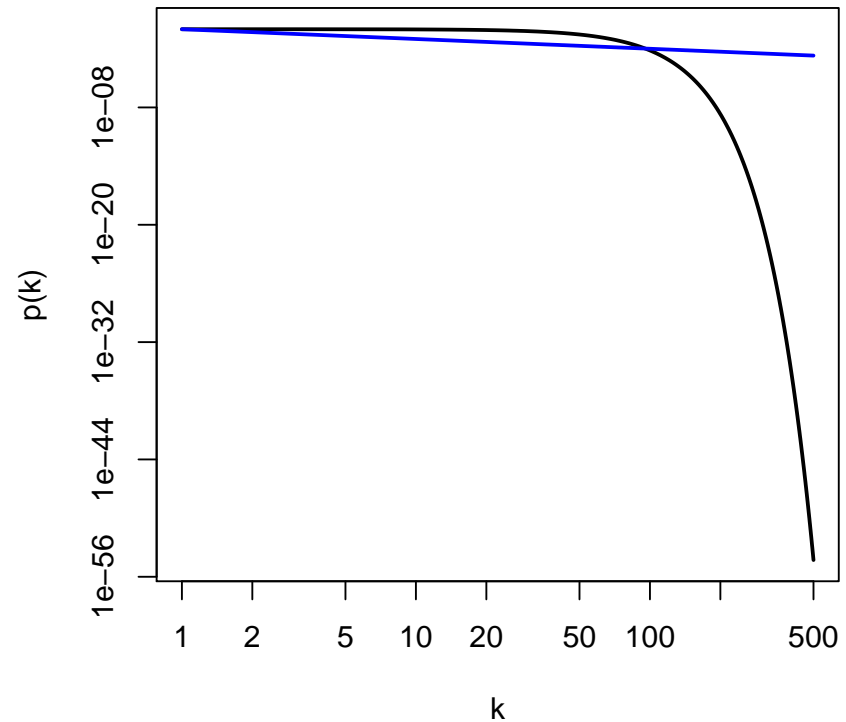
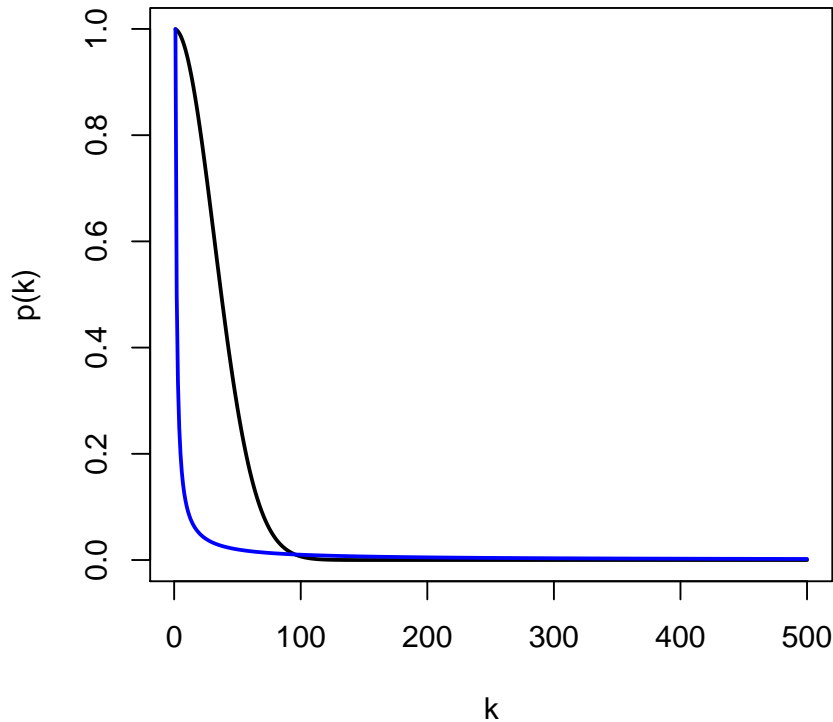
$$p_k \sim k^{-\gamma}$$

$$\ln p_k \sim -\gamma \ln k$$



Power Laws versus Bell Curves: “Heavy tails”

- Power law distribution: $p_k \sim k^{-\gamma}$.
- Gaussian distribution: $p_k \sim \exp(-k^2/2\sigma^2)$.



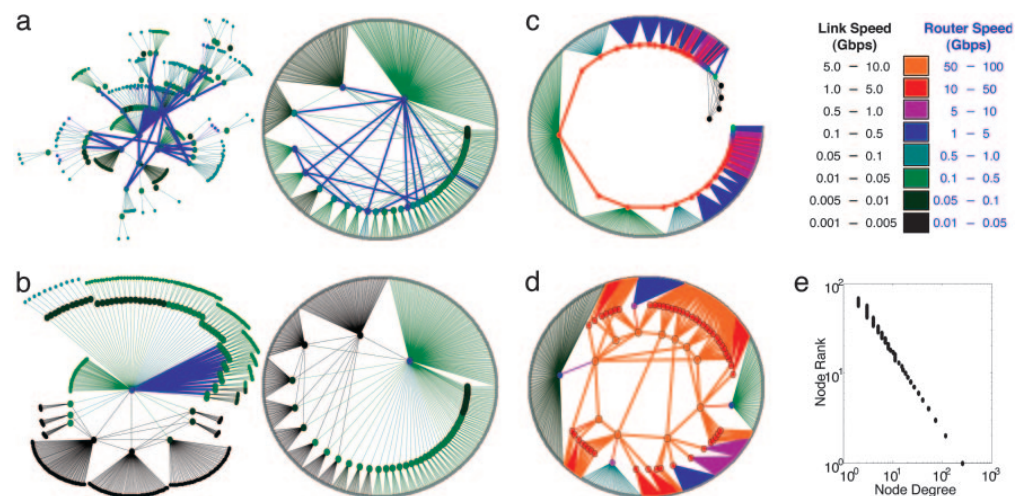
- Most nodes have low degree
- But a few nodes are hubs, with massive degree

Many network growth models produce power law degree distribution (we will study some of these)

- Preferential attachment
- Copying models (WWW, biological networks, ...)
- Optimization models

Degree distribution misses other structure.

- Doyle, et. al.,
PNAS **102** (4)2005.



Power law probability distributions: $p_k = Ak^{-\gamma}$ with $\gamma > 0$

- $0 \leq p_k \leq 1 \quad \forall k$ which are valid degrees (typically $k \in \mathbb{Z}^+$).
- Must be properly normalized:

$$\sum_{k=1}^{\infty} p_k = \sum_{k=1}^{\infty} \frac{A}{k^{\gamma}} = 1$$

- Approximating discrete sum by integral:

$$\begin{aligned} 1 &= \int_{k=1}^{\infty} \frac{A}{k^{\gamma}} = - \left(\frac{A}{\gamma - 1} \right) \frac{1}{k^{(\gamma-1)}} \Big|_{k=1}^{\infty} \\ &= \left(\frac{A}{1 - \gamma} \right) \left(\frac{1}{\infty^{(\gamma-1)}} - 1 \right) \end{aligned}$$

- Finite requirement means $\gamma > 1$, in which case $A = (\gamma - 1)$.

The first moment (the mean)

Recall, $p_k = \frac{A}{k^\gamma}$

- Mean degree:

$$\langle k \rangle = \sum_{k=1}^{\infty} k p_k \approx \int_{k=1}^{\infty} k p_k dk = \int_{k=1}^{\infty} \frac{A}{k^{(\gamma-1)}} dk$$

Diverges (i.e., $\langle k \rangle \rightarrow \infty$) if $\gamma \leq 2$.

The second moment and the variance

- Second moment:

$$\langle k^2 \rangle = \sum_{k=1}^{\infty} k^2 p_k \approx \int_{k=1}^{\infty} k^2 p_k dk = \int_{k=1}^{\infty} \frac{A}{k^{(\gamma-2)}} dk$$

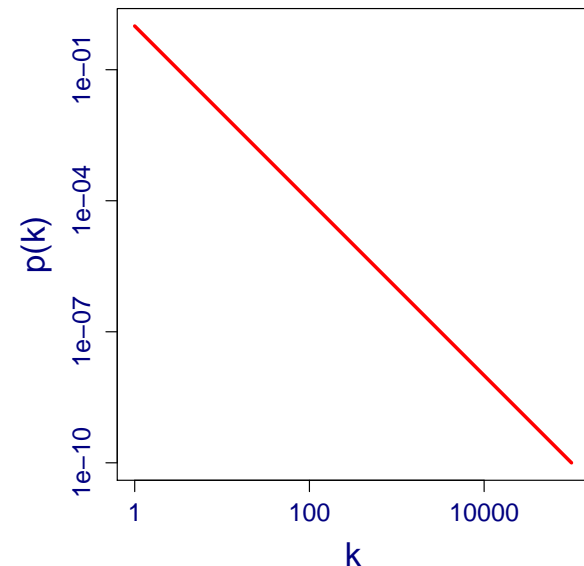
Diverges (i.e., $\langle k^2 \rangle \rightarrow \infty$) if $\gamma \leq 3$.

- Variance = $\langle k^2 \rangle - \langle k \rangle^2$, likewise diverges if $\gamma \leq 3$.

Properties of a power law PDF (Summary)

(PDF = probability density function)

- To be a properly defined probability distribution need $\gamma > 1$.
- For $1 < \gamma \leq 2$, both the average $\langle k \rangle$ and variance σ^2 are infinite!
- For $2 < \gamma \leq 3$, average $\langle k \rangle$ is finite, but variance σ^2 is infinite!
- For $\gamma > 3$, both average and variance finite.



Why a power law is “scale-free”

- Power law for “ x ”, means “scale-free” in x :

$$p(bx) = (bx)^{-\gamma} = b^{-\gamma}p(x)$$

$$\boxed{\frac{p(bk)}{p(k)} = b^{-\gamma}} \text{ regardless of } k.$$

In contrast consider: $p(k) = A \exp(-k)$.

So $p(bk) = A \exp(-bk)$.

$$\boxed{\frac{p(bk)}{p(k)} = \exp[-k(b-1)]} \text{ dependent on } k$$

Self-similar/scale-free fractal structures



Sierpinski Sieve/Gasket/Fractal, $N \sim r^d$.

When r doubles, N triples: $3 = 2^d$

$$d = \log N / \log r = \log 3 / \log 2$$

Power laws in the real world

Confusion

- Power law
- Log normal
- Weibull
- Stretched exponential

All of these distributions can look the same! (Especially when we are dealing with finite data sets — not enough data to get good statistics).

How to deal with real data

- Can adjust bin size: increase exponentially with degree.
- Consider the Cumulative PDF (the CDF): $P_k = \sum_{l=k}^{\infty} p_l$.

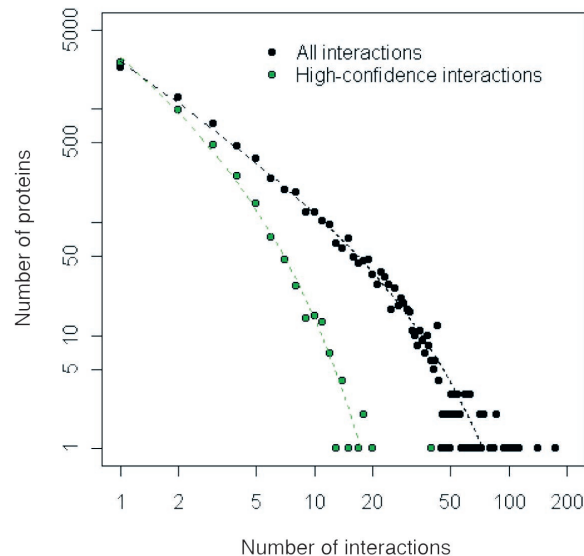
Good reviews:

- Aaron Clauset, Cosma R. Shalizi, M. E. J. Newman. “Power-Law Distributions in Empirical Data”, *SIAM Review*, Vol. 51, No. 4. (2009), pp. 661-703.
- A Brief History of Generative Models for Power Law and Lognormal Distributions Michael Mitzenmacher, *Internet Math*. Vol 1 (2003), 226-251.

Power law with exponential tail

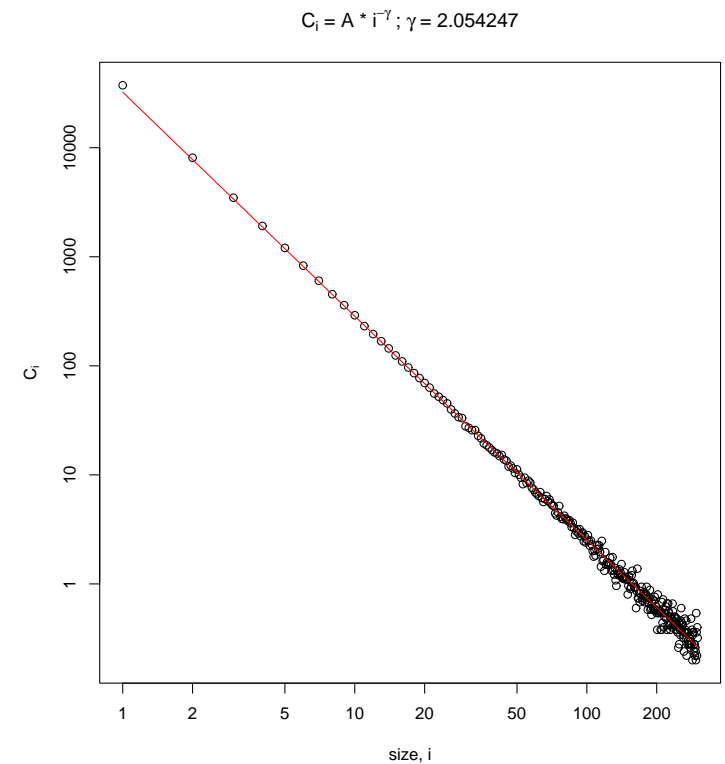
Ubiquitous empirical measurements:

System with: $p(x) \sim x^{-B} \exp(-x/C)$	B	C
Full protein-interaction map of <i>Drosophila</i>	1.20	0.038
High-confidence protein-interaction map of <i>Drosophila</i>	1.26	0.27
Gene-flow/hybridization network of plants as function of spatial distance	0.75	10^5 m
Earthquake magnitude	1.35 - 1.7	$\sim 10^{21}$ Nm
Avalanche size of ferromagnetic materials	1.2 - 1.4	$L^{1.4}$
ArXiv co-author network	1.3	53
MEDLINE co-author network	2.1	~ 5800
PNAS paper citation network	0.49	4.21



True power laws are observed in many systems

- Signature of a system at the “critical point” of a phase transition.
Foundation of renormalization group approach to critical phenomena
- Random graphs at critical point;
component sizes: $N_k \sim k^{-5/2}$
(Note, $\gamma = 2.5$)



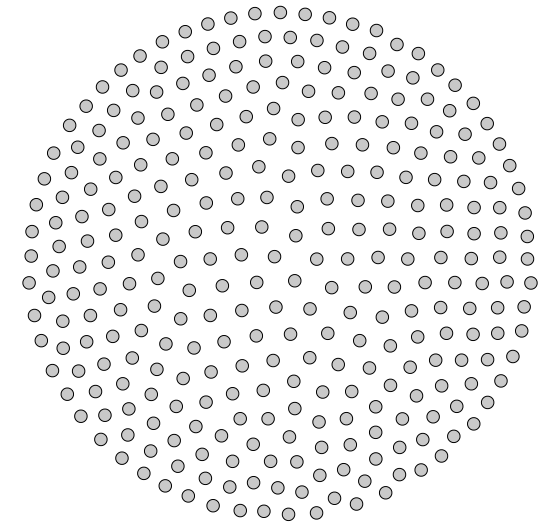
The origins of network theory: Random graphs

What does a “typical” graph with n vertices and m edges look like?

- P. Erdős and A. Rényi, “On random graphs”, *Publ. Math. Debrecen.* **6**, 1959.
- P. Erdős and A. Rényi, “On the evolution of random graphs”, *Publ. Math. Inst. Hungar. Acad. Sci.* **5**, 1960.
- E. N. Gilbert, “Random graphs”, *Annals of Mathematical Statistics* **30**, 1959.

Erdős-Rényi random graphs

- Consider a *labelled* graph. Each vertex has a label ranging from $[1, 2, 3, \dots, n]$, for a set of n vertices. (This will make counting and analysis easier.)



- Let E denote the total number of edges possible:

$$E = \binom{N}{2} = \frac{N!}{2!(N-2)!} = \frac{N(N-1)}{2}$$

(If directed edges, we would not divide by 2).

Two formulations

- 1) $\mathcal{G}(n, p)$: The *ensemble* of graphs constructed by putting in edges with probability p , independent of one another. (An edge is present with probability p and absent with probability $[1 - p]$.)

Let $G(n, p)$ denote a random realization of $\mathcal{G}(n, p)$.

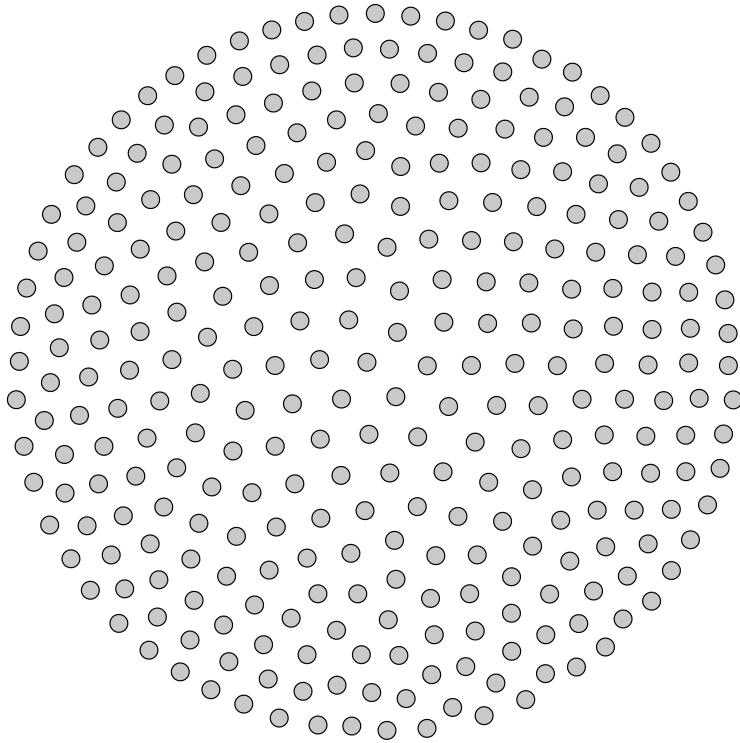
- 2) $\mathcal{G}(n, m)$: The ensemble of all graphs with n nodes and exactly m edges.

Let $G(n, m)$ denote a random realization of $\mathcal{G}(n, m)$.

- The two are almost interchangeable with $m = pE$. (Recall, E is total number of edges possible).
- We will focus on $G(n, p)$.

The “classic” random graph, $G(N, p)$ (The Null Model)

- P. Erdős and A. Rényi, “On random graphs”, *Publ. Math. Debrecen*. 1959.
- P. Erdős and A. Rényi, “On the evolution of random graphs”, *Publ. Math. Inst. Hungar. Acad. Sci.* 1960.
- E. N. Gilbert, “Random graphs”, *Annals of Mathematical Statistics*, 1959.



- Start with N isolated vertices.
- Add random edges one-at-a-time.
 $N(N - 1)/2$ total edges possible.
- After E edges, probability p of any edge is $p = 2E/N(N - 1)$

What does the resulting graph look like?

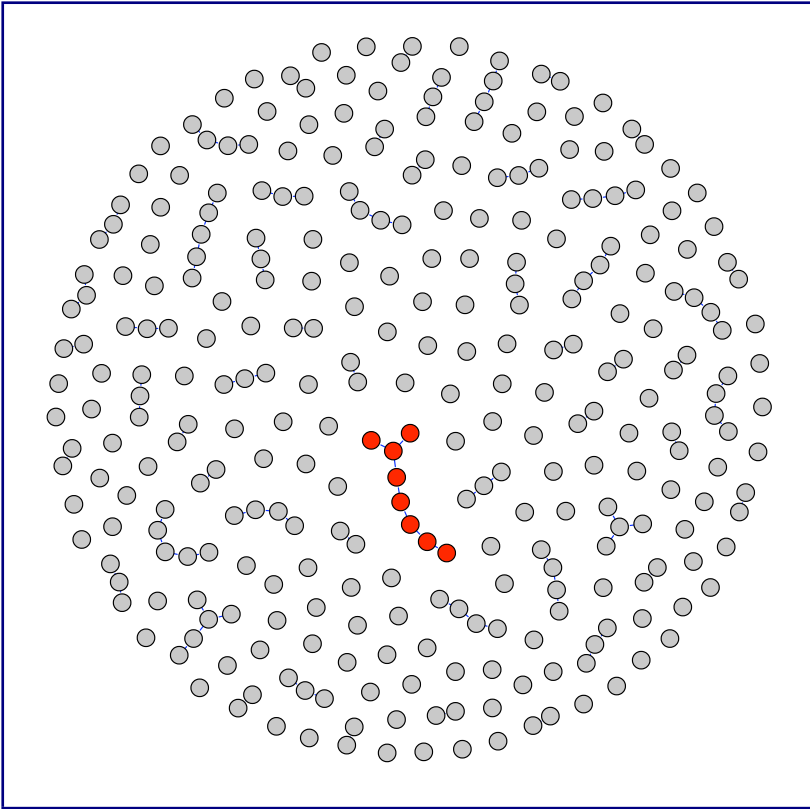
(Typical member of the ensemble)

Explicitly building $G(n, p)$

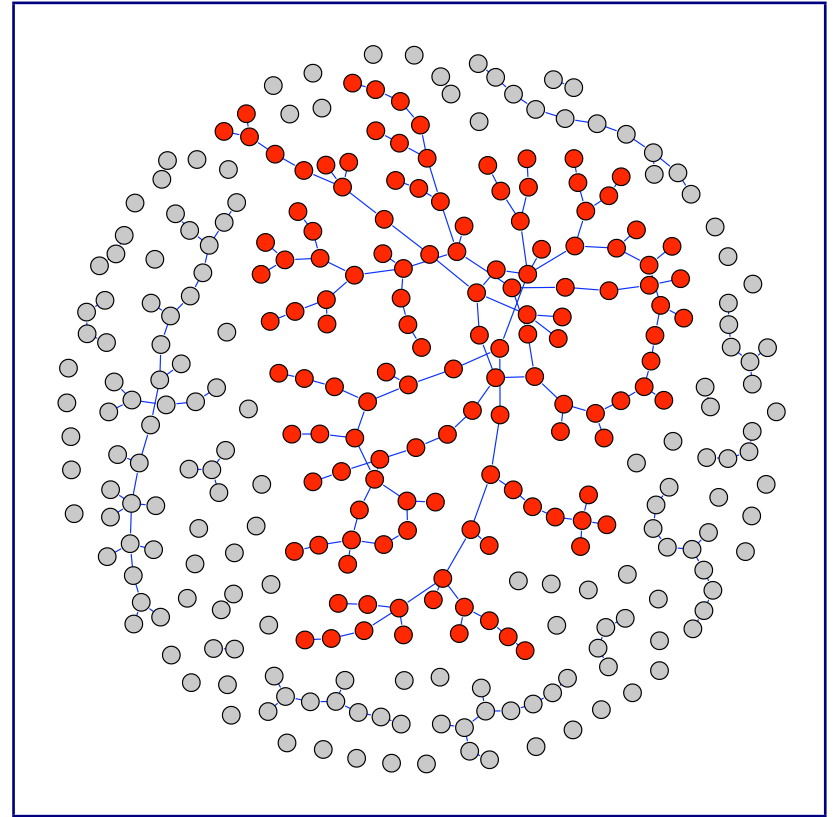
- Build a realization of $G(n, p)$ by the following graph process:
- Start with n isolated vertices.
- At each discrete time step, add one edge chosen at random from edges not yet present on the graph.
- At “time” t (i.e., at the addition of t edges), we have built a realization of $G(n, p)$ where $p = t/E$.
- This is a Markov process (build graph at time $t + 1$ from graph at time t).

Ben-Naim, **Krapivsky**, “Kinetic theory of random graphs”, *PRE*,
2005.

N=300



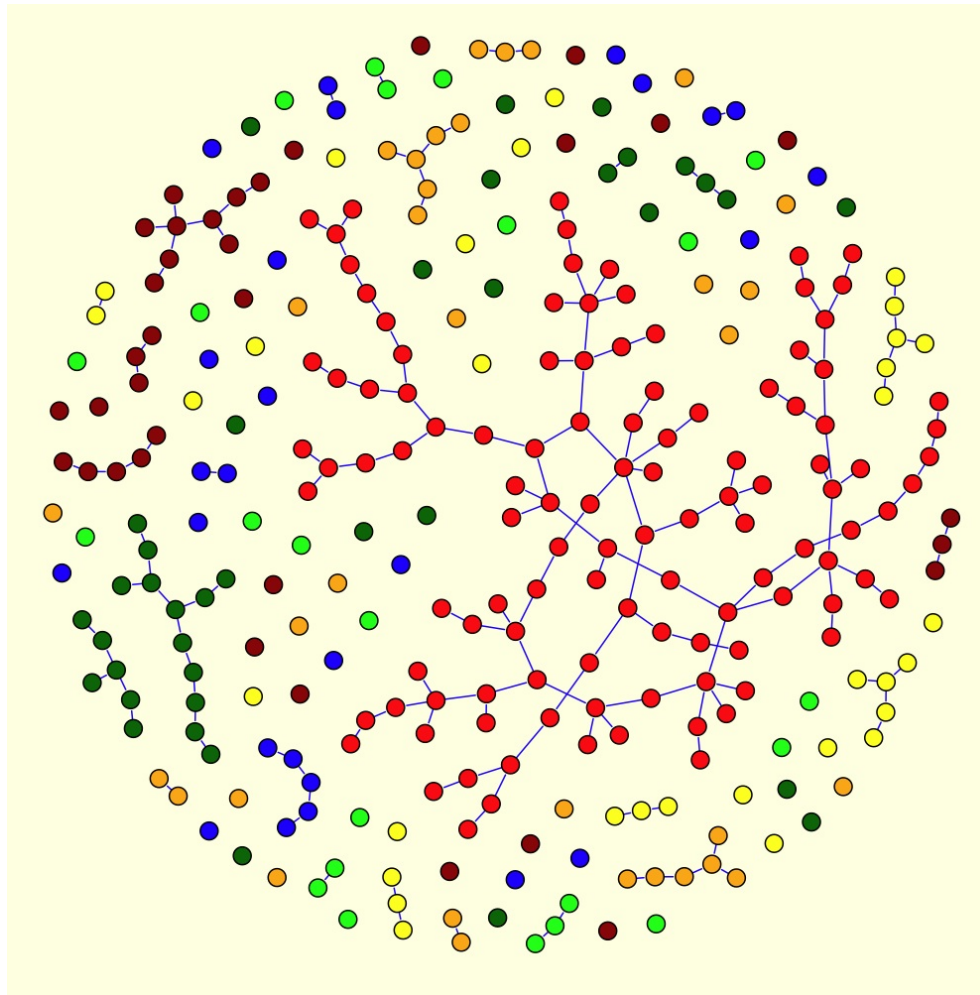
$$p = 1/400 = 0.0025$$



$$p = 1/200 = 0.005$$

Component

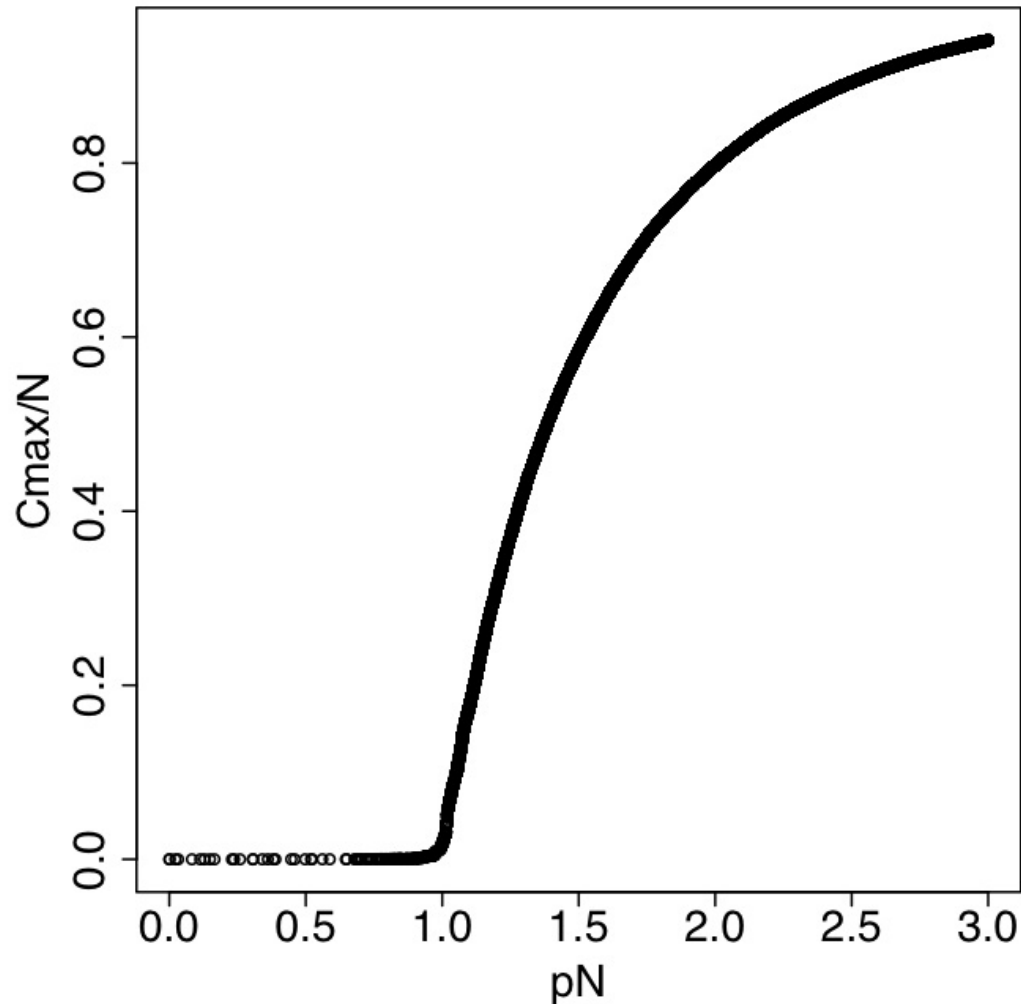
A **component** is a subset of vertices in the graph each of which is reachable from the other by some path through the network.



Behavior for small p

- Consider a realization $G(n, p)$ for $0 < p < 1$ and $n \rightarrow \infty$.
(A number of interesting properties of random graphs can be proven in this limit.
The $n \rightarrow \infty$ limit is also called the “thermodynamic limit”.)
- Let $C_{max}(p)$ denote the size of the largest component of $G(n, p)$ as a function of p .
- For small p , few edges on the graph. Almost all vertices disconnected. The components are small, with size $O(\log n)$, independent of p .
- Keep increasing p (or equivalently t in our model).
At $p = 1/n$ (i.e. $t = E/n$), something surprising happens:

Emergence of a “giant component”



- $p_c = 1/N$.
- $p < p_c$, $C_{\max} \sim \log(N)$
- $p > p_c$, $C_{\max} \sim A \cdot N$

(Ave node degree $t = pN$
so $t_c = 1$.)

Branching process (Galton-Watson); “tree”-like at $t_c = 1$.

A Phase Transition!

An abrupt sudden change in one or more physical properties, resulting from a small change in an external control parameter.

Examples from physical systems:

- Magnetization
- Superconductivity
- Liquid/Gas
- Bose-Einstein condensation

Giant component observed in real-world networks

- Formation reminiscent of many real-world networks.
“Gain critical mass”.
- Lower bound on emergence of epidemic outbreak.
- The giant component/Strongly Connected Component used extensively to categorize networks.

Phase transition in connectivity

- Below $p = 1/n$, only small disconnected components.
- Above $p = 1/n$, one large component, which quickly gains more mass. All other components remain sub-linear.
- Note the average node degree, z :

$$\begin{aligned} z &= (2 \times \#edges) / \#vertices \\ &= (2pn(n-1)/2) / n = pn(n-1)/n = (n-1)p \approx np. \end{aligned}$$

(Factor of 2 since each edge contributes degree to two vertices – each end of the edge contributes.)

Recall, expected number of edges, is $pn(n-1)/2$.

- At the phase transition, $z = np = 1$. The phase transition occurs when the average vertex degree is one!

Erdős-Rényi, a continuous, second order transition: Mean-field scaling behaviors

- Divergence of susceptibility: $\chi = \frac{\partial m}{\partial h} \sim |T - T_c|^{-\gamma}$
- Random graph “susceptibility” (second moment of the component sizes): $\chi = \sum_{i=1}^{\infty} i^2 n_i$
- For Erdős-Rényi, $\chi \sim |t_c - t|^{-\gamma}$, with $\gamma = 1$.

Power law correlation lengths and response functions →

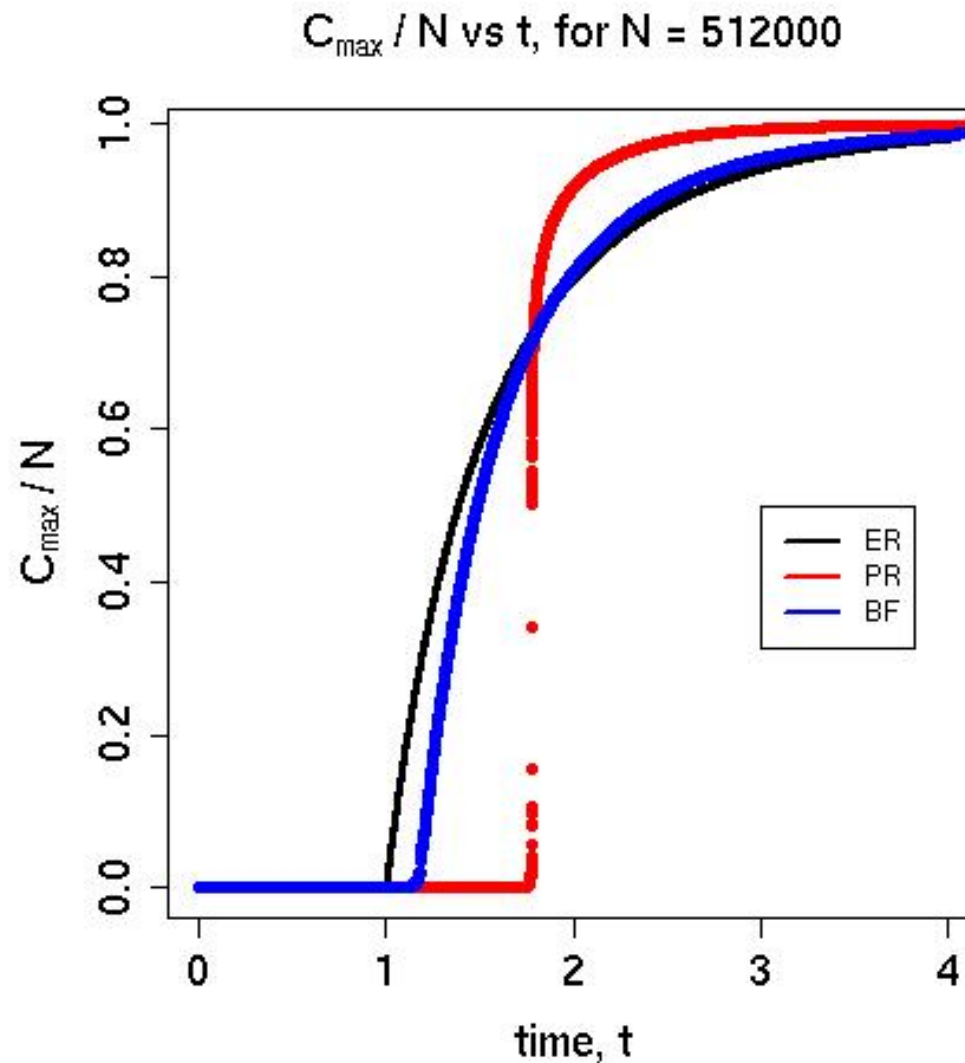
Potential EARLY WARNING SIGNALS

(e.g., Scheffer et al. *Nature* 461, 2009)

Is connectivity a good thing?

- Communication, transportation networks
- Spreading of a virus (human or computer)

Algorithms for suppressing the emergence of the Giant Component

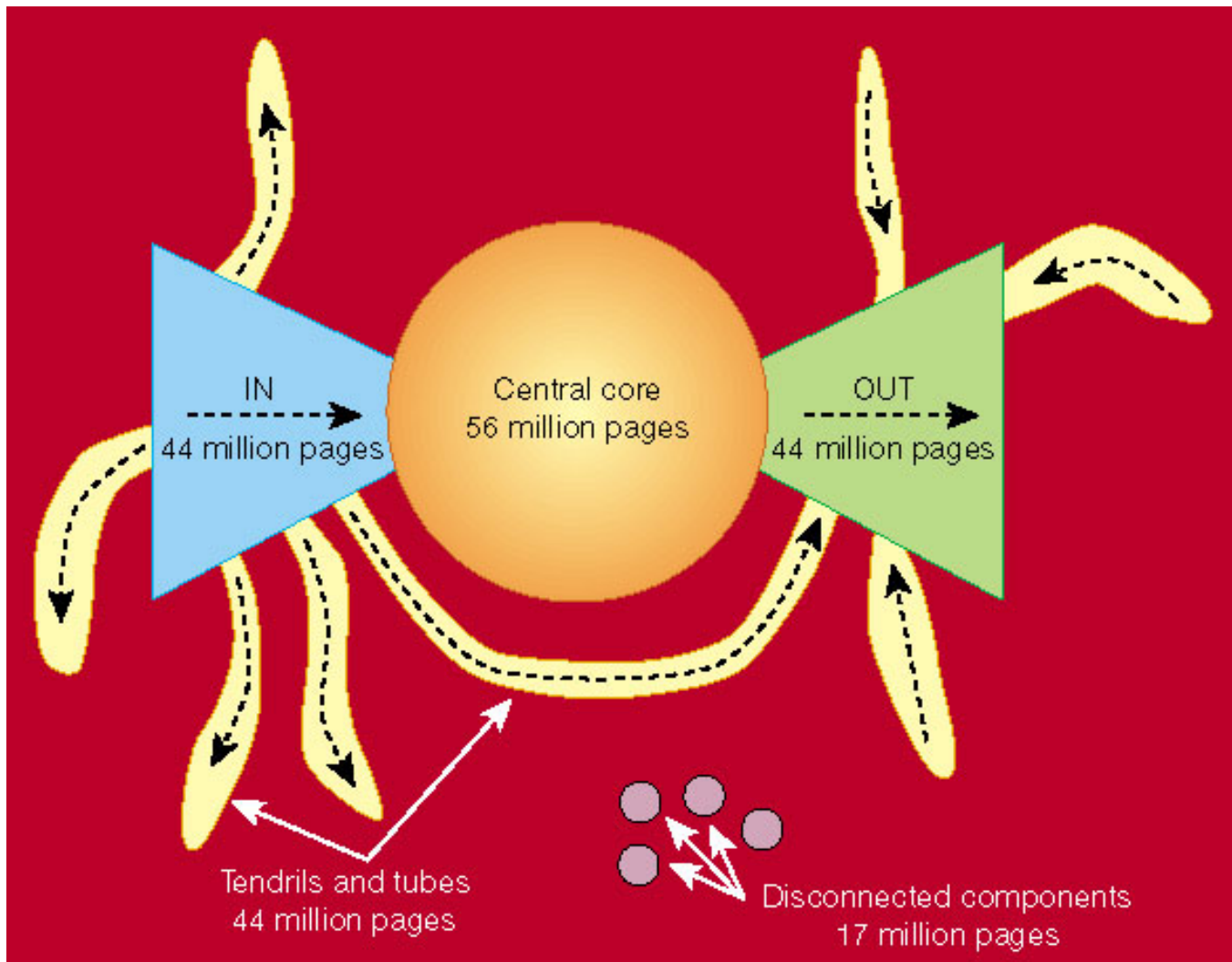


e.g. “Explosive percolation”, Achlioptas, D’Souza, Spencer,
Science, 2009.

Random graphs as real-world networks?

- What about degree distribution, clustering, assortativity....?
 - Shown later, Erdos-Renyi yields a Poisson degree distribution, but “configuration” models work around this.
 - Still need null models to match other properties.
- e.g., “Network Analysis in the Social Sciences”, S. P. Borgatti, A. Mehra, D. J. Brass, G. Labianca, *Science* **323**, 892-895, 2009.
 - Why would a real network look like a random one?
 - Local properties of nodes and edges, not statistics of the network.
- Developing the correct null models?

The giant component/Strongly Connected Component of the WWW



From "The web is a bow tie" Nature **405**, 113 (11 May 2000)

Summary: Terms introduced today

- Component
- Phase transition
- Degree distribution
- Graph diameter

Further reading on random graphs

- M. E. J. Newman review, pages 20-25. (Heuristic arguments)
- R. Durrett book, Chaps 1 and 2. (Technical proofs)
- B. Bollobás, *Random Graphs*, 2nd Edition, Cambridge U Press, 2001 (the seminal text on the mathematics of random graphs).

Class structure

- Two tracks to the class:

Track A: Project

- (1) Common homeworks (e.g. HW1.pdf, HW2.pdf) and
- (2) HW1a.pdf, HW2a.pdf etc.

Track B: Advanced HWs

- (1) Common homeworks (e.g. HW1.pdf, HW2.pdf) and
- (2) HW1b.pdf, HW2b.pdf etc.

- Track A: Project

- Teams of 5-6 people ideal
- Negative results are OK
- Ideally aim to have a result for a journal or conference

e.g., “Latent social structure in open source projects”, C Bird, D Pattison, R D’Souza, V Filkov, P Devanbu, ACM SIGSOFT 2008.

Project pitch – HW1a

- One page describing your idea. Submitted via Canvas and shared with the class.
- Skill sets to merge:
Domain specific questions / Methods / Data sets

Jumpstart: In-class on Monday (April 10th) — pitch your idea and build a team!

Possible topic areas

- Transportation networks and flows; multi-modal transportation
- Open source software – e.g., social and technological networks in github
- Machine learning – e.g., bring network connectivity into binary classifiers
- Power grid modeling
- Opinion dynamics / social unrest / multiplex opinion dynamics
- Ranking in networks; especially temporal, multilayered, higher-order
- Multilayered and temporal macaque monkey networks
- Shocks and tipping points
- Extend standard metrics to multilayered, temporal, or higher-order networks
- Co-author and citation networks
- Food networks
- Recommendation systems
- Biological networks
- Terrorist networks
- See also class homepage “Projects” tab