

## Quiz 1

First name:  Last name:  Seat #:

Instructions: No notes/books/gadgets/neighbors.

1. Prof. Rogaway criticized the book's definition of a *language* as a *set of strings*; he said that one should add in that .
2. Let  $L = \{10,001\}$ . List the first six strings of  $L^*$  in lexicographic order (assume  $0 < 1$ ):
3. A **DFA** is a five-tuple  $M = (Q, \Sigma, \delta, q_0, F)$  where  $Q$  is a finite set,  $\Sigma$  is an alphabet,  $q_0 \in Q$ ,  $F \subseteq Q$ , and  $\delta$  is a function with domain  and range . In contrast, the transition function for an **NFA** has domain  and range .
4. **Left side:** draw a **DFA** for the language  $L = \{0,1\}^* \{111\}$ . Make your DFA as small (= as few states) as possible. **Right side:** Draw an **NFA** for  $L$ . Make your NFA as small (= as few states) as possible. Among all such NFAs, use as few transitions as possible.

5. **Darken** the **correct** box. No justification is required. If you're not sure, guess.

- |     |                               |                                |   |
|-----|-------------------------------|--------------------------------|---|
| (a) | <input type="checkbox"/> True | <input type="checkbox"/> False | There's a language $L$ such that $L = L^R$ (the reversal of $L$ , $\{x^R : x \in L\}$ )   |
| (b) | <input type="checkbox"/> True | <input type="checkbox"/> False | There's a language $L_0$ that's a subset of every language.   |
| (c) | <input type="checkbox"/> True | <input type="checkbox"/> False | If $L$ is accepted by a DFA then $\bar{L}$ is accepted by a DFA.  |
| (d) | <input type="checkbox"/> True | <input type="checkbox"/> False | If $ L  = 5$ then $L^*$ is infinite.  |
| (e) | <input type="checkbox"/> True | <input type="checkbox"/> False | The concatenation of finite languages $A$ and $B$ is finite.  |
| (f) | <input type="checkbox"/> True | <input type="checkbox"/> False | Every finite language $L$ is accepted by some NFA.  |
| (g) | <input type="checkbox"/> True | <input type="checkbox"/> False | There's an algorithm to determine if a multivariate polynomial has an integer root, but the best known algorithm is extremely slow. |
| (h) | <input type="checkbox"/> True | <input type="checkbox"/> False | If $M = (Q, \Sigma, \delta, q_0, F)$ is a DFA and $F = Q$ then $L(M) = \Sigma^*$ .  |
| (i) | <input type="checkbox"/> True | <input type="checkbox"/> False | If $M = (Q, \Sigma, \delta, q_0, F)$ is an NFA and $F = Q$ then $L(M) = \Sigma^*$ .   |
| (j) | <input type="checkbox"/> True | <input type="checkbox"/> False | If $\Sigma$ is an alphabet then $\Sigma^* - \Sigma^+ = \{\varepsilon\}$ .   |

6. In PS #2 we defined  $\mathcal{E}(L) = \{x \in L : \text{there exists a } y \in \Sigma^+ \text{ for which } xy \in L\}$ .

Let  $L = \{0\}^* \{1\}$ . Then  $\mathcal{E}(L) =$  . (Write it in a simple form.)