## Quiz 4

Firstname Lastname: \_\_\_\_\_ ID# \_\_\_\_\_ Seat# \_\_\_\_

- Don't sit next to anyone you know.

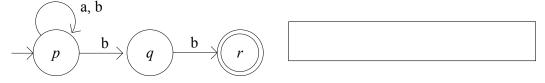
- Don't turn over this paper until you are asked to.

- When you finish, put this side up once again.

- Most or all problems will be graded all-or-northing.

- Relax, these quizzes are too insignificant to get stressed over.

Happy Friday! phil rogaway (1) Write a regular expression for the language of the following NFA  $M_1$ . Make it as simple as possible. Use standard abbreviations, not writing the concatenation symbol or extra parentheses.



(3) Darken the box to indicate if the statement is True or False. Really make your mark dark. As always, a statement is True if it is always true; otherwise it is False.

True False Every regular language can be accepted by a DFA with an odd number of states.

True False Every regular language can be accepted by a DFA whose start state is never visited twice.

(4) Same instructions. Throughout, fix an NFA  $M = (Q, \Sigma, \delta, q_0, F)$ .

True False Suppose there is a  $q_0 \rightsquigarrow q$  path in the diagram for M where  $q \in F$  and the concatenation of edge-labels along the path is s. Then M accepts s.

<u>True</u> False Suppose there is a  $q_0 \rightsquigarrow q$  path in the diagram for M where  $q \notin F$  and the concatenation of edge-labels along the path is s. Then M rejects s.

(5) Let's begin a proof that  $L = \{a^n b^n : n \ge 0\}$  is **not** regular: Assume for contradiction that L is regular. Then there is a DFA  $M = (Q, \Sigma, \delta, q_0, F)$  that accepts L. Let N = |Q|. Consider the N + 1 strings  $\square$ . Each of these strings w determines a state  $\delta^*(q_0, w)$ . By the pigeonhole principle (PHP), we know that some two of these states  $\square$ . And so on ...

Make sure that whatever you write gives grammatical English.