## Quiz 4

Firstname Lastname:
ID\#
Seat\# -

- Don't sit next to anyone you know.
- Don't turn over this paper until you are asked to.
- When you finish, put this side up once again.
- Most or all problems will be graded all-or-northing.
- Relax, these quizzes are too insignificant to get stressed over.

Happy Friday!
phil rogaway
(1) Write a regular expression for the language of the following NFA $M_{1}$. Make it as simple as possible. Use standard abbreviations, not writing the concatenation symbol or extra parentheses.

$\square$
(2) Suppose $L$ is accepted by a 10 -state DFA $M_{2}$. Using constructions described in class, we could convert $M_{2}$ into a $\square$-state NFA $M_{2}^{\prime}$ for $L \bar{L}$ (the bar denoting the complement of the language). We could then convert $M_{2}^{\prime}$ into a $\square$-state DFA for $L \bar{L}$.
(3) Darken the box to indicate if the statement is True or False. Really make your mark dark. As always, a statement is True if it is always true; otherwise it is False.

True False Every regular language can be accepted by a DFA with an odd number of states. | True | False | Every regular language can be accepted by a DFA whose start state is never |
| :--- | :--- | :--- | visited twice.

(4) Same instructions. Throughout, fix an NFA $M=\left(Q, \Sigma, \delta, q_{0}, F\right)$.

True False Suppose there is a $q_{0} \rightsquigarrow q$ path in the diagram for $M$ where $q \in F$ and the concatenation of edge-labels along the path is $s$. Then $M$ accepts $s$.
True False Suppose there is a $q_{0} \rightsquigarrow q$ path in the diagram for $M$ where $q \notin F$ and the concatenation of edge-labels along the path is $s$. Then $M$ rejects $s$.
(5) Let's begin a proof that $L=\left\{a^{n} b^{n}: n \geq 0\right\}$ is not regular:

Assume for contradiction that $L$ is regular. Then there is a DFA $M=\left(Q, \Sigma, \delta, q_{0}, F\right)$ that accepts $L$. Let $N=|Q|$. Consider the $N+1$ strings $\square$ Each of these strings $w$ determines a state $\delta^{*}\left(q_{0}, w\right)$. By the pigeonhole principle (PHP), we know that some two of these states $\square$. And so on ...
Make sure that whatever you write gives grammatical English.

