Quiz 2 tomorrow & PS4

- Q2: review PS1-3 and quiz 1, especially difficult questions, definitions from class
- PS4 Q5: calculus may be used; Q8c ⊊ **proper subset** of, *A* ⊊ *B* means that *A* is a subset of but not equal to *B*

## Mathematical Induction

To prove proposition P(n) for all integers  $n \ge n_0$ , prove the <u>basis</u>:  $n = n_0$  is true <u>Inductive step</u>: suppose n = k is true, show n = k + 1 is true

1. Prove that  $n^3 + 5n$  is divisible by 6 for every  $n \ge 1$ .

<u>Basis</u>: n = 1:  $6 | 1^3 + 5(1)$ <u>Inductive step</u>: suppose n = k is true, so  $6 | k^3 + 5k$ Check n = k + 1:  $(k + 1)^3 + 5(k + 1) = (k^3 + 3k^2 + 3k + 1) + (5k + 5)$   $= (k^3 + 5k) + (3k^2 + 3k + 6)$   $= (k^3 + 5k) + 3k(k + 1) + 6$ From the inductive hypothesis,  $6 | (k^3 + 5k) \cdot k(k + 1)$  is even, so 6 | 3k(k + 1),

and 6 | 6. Inductive step verified.

2. Prove that  $8^n - 3^n$  is divisible by 5 for every  $n \ge 1$ .

Basis:  $n = 1:5 | 8^{1} - 3^{1}$ Inductive step: suppose n = k is true, so  $5 | 8^{k} - 3^{k}$ Check n = k + 1:  $8^{k+1} - 3^{k+1} = 8 * 8^{k} - 3 * 3^{k}$   $= 8 * 8^{k} - 3 * 8^{k} + 3 * 8^{k} - 3 * 3^{k}$   $= (8 - 3) * 8^{k} + 3 * (8^{k} - 3^{k})$  $= 5 * 8^{k} + 3 * (8^{k} - 3^{k})$ 

We know that  $5 | 5 * 8^k$ . From the inductive hypothesis,  $5 | (8^k - 3^k)$ . Inductive step verified.

3. Prove that  $n^2 > 7n + 1$  for all  $n \ge 8$ .

<u>Basis</u>: n = 8: 64 > 56 + 1<u>Inductive step</u>: suppose n = k is true, so  $k^2 > 7k + 1$ Check n = k + 1,  $(k + 1)^2 > 7(k + 1) + 1$ LHS:  $(k + 1)^2 = k^2 + 2k + 1 > (7k + 1) + 2k + 1 = 7(k + 1) + 1 + (2k - 6)$ Since  $k \ge 8$ , 2k - 6 > 0 so LHS is greater than RHS, inductive step verified.

4. Induction on a definition An expression E is defined as

 $E \rightarrow < Number >$  $E \rightarrow (E + E) \mid (E * E) \mid (E^{E}) \mid \ln(E)$  Prove that an expression *E* has the same number of (and).

<u>Basis</u>: E = < Number >, 0 ('s and )'s

<u>Inductive step</u>: suppose n = k is true,  $E_k$  has the same number of ( and ). Check n = k + 1,  $E_{k+1}$ :

Take on either the form  $(E + E) | (E * E) | (E^{E}) | \ln(E)$ , all of which contain an equal number of ( and ). Inductive step verified.

Strong vs. Weak Induction

- method so far is sometimes called **weak induction**,  $P(k) \rightarrow P(k+1)$
- **strong induction** shows  $(\forall n)P(n)$  by showing  $P(1) \land ... \land P(k) \rightarrow P(k+1)$
- both approaches are equally "strong", strong induction is easier for certain proofs, such as PS4 Q6
- 5. Prove that  $a_n \leq 2^n$  for the sequence  $a_0 = 1$ ,  $a_1 = 2$ ,  $a_2 = 3$ ,  $a_i = a_i + a_{i-1} + a_{i-2}$ <u>Basis</u>: n = 0,1,2 true because  $1 \leq 2^0, 2 \leq 2^1, 3 \leq 2^2$ <u>Inductive step</u>: assume  $n = 0 \dots (k - 1)$  is true, show that n = k is true i.e. show that  $P(k) = a_k \leq 2^k$

$$\begin{aligned} a_k &= a_{k-1} + a_{k-2} + a_{k-3} \\ &\leq 2^{k-1} + 2^{k-2} + 2^{k-3} \\ &\leq 2^0 + 2^1 + \dots + 2^{k-3} + 2^{k-2} + 2^{k-1} \end{aligned}$$
  
Geometric series:  $\leq 2^k - 1 \leq 2^k$ 

## Envelope substitution question

A store sells envelopes in packages of 5 and 12. Prove that for  $\ge n$ , the store can sell exactly *n* envelopes. To find *n* without trying every number, start with the smallest substitutions to make one more envelope (=inductive step).

The lowest substitution going from 5- to 12-packs is  $7(5) \rightarrow 3(12)$ . This substitution needs at least 7 5-packs.

Consider the case with less than 7 5-packs, i.e. at most 6 5-packs ( $\leq$  30). The lowest substitution going from 12- to 5-packs is 2(12)  $\rightarrow$  5(5). This substitution needs at least 2 12-packs ( $\geq$  24).

To find *n*, we need only to check  $30 < n \le 54$ , since there will be numbers below 30 that cannot be achieved, and any number 54 or above will work for certain. We will find n = 44 to be the first number beyond which all *n* is satisfied.

We can also reverse the order of substitutions. In this example that would be checking  $12 < n \le 47$ , which also includes n = 44.

Once n is found, we can formally prove our hypothesis by induction using n as the basis.